

Uni-critical Branner-Hubbard conjecture

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1 Introduction

The Branner-Hubbard conjecture has been recently proved in its full generality (see [KS], [QY]):

Theorem (Qiu-Yin, Kozlovski-van Strien). *Let $f : \mathbb{C} \rightarrow \mathbb{C}$ be a complex polynomial. Denote by K_f its filled Julia set. Assume that every component of K_f containing a critical point is aperiodic under the iteration of f . Then K_f is a Cantor set.*

A particular case of this result is the uni-critical one:

Theorem 1.1 *Assume that $f : \mathbb{C} \rightarrow \mathbb{C}$ is a polynomial with all but one critical point c escaping to ∞ . The non-escaping critical point c might have multiplicity. Assume that the K_f -component of c is not periodic. Then K_f is a Cantor set.*

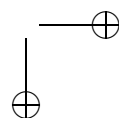
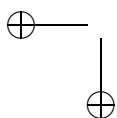
The case c is a simple critical point was originally proved by Branner-Hubbard ([BH]), which lead them to the conjecture for a general polynomial.

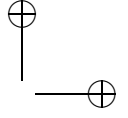
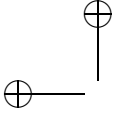
In this short note we will illustrate the new techniques involved in the proof of this conjecture, by proving a key statement (Theorem 1.2 below) that leads easily to the unicritical Theorem 1.1. The proof we present here follows essentially the same line as [QY].

These new techniques, namely Kozlovski-Shen-Strien nest and Kahn-Lyubich covering lemma, when combined with various other techniques, have lead to several interesting results in the field. For a detailed account, see for example [QY].

2 Results

We set up a puzzle by taking an equipotential of f so that the bounded pieces contain no other critical points than c . These pieces consist of our puzzle of depth 0. A connected component under f^{-n} is a puzzle piece of depth n .





The tableau $T(c)$ is by definition a grill indexed by $-\mathbb{N} \times \mathbb{N}$, where $\mathbb{N} = \{0, 1, \dots\}$, with the 0-th column representing the nest of the consecutive puzzle pieces containing the critical point c . Now let I be a critical puzzle piece of depth, say, m . It occupies therefore the $(-m, 0)$ th entry of $T(c)$. We use $f^j(I)$, $j = 1, \dots, m$ to form the north-east diagonal segment starting from I , so that $f^j(I)$ occupies the $(-m + j, j)$ th entry of $T(c)$.

A position in $T(c)$ is marked by \circ if it is non-critical, by \bullet if it is a critical puzzle piece, and by \times if unknown.

A *child* of a critical puzzle piece I is defined to be a critical puzzle piece J with deeper depth so that, setting $k = |I - J|$, $f^j(J)$, $j = 1, 2, \dots, k - 1$, is not a critical piece, and $f^k(J) = I$. We will use the symbol $J \xrightarrow{f^k} I$. In other words $f : J \rightarrow f(J)$ is non-univalent but $f^{k-1} : f(J) \rightarrow I$ is univalent. We will use δ to denote the degree of f . Then $\deg(J \xrightarrow{f^k} I) = \delta$.

Viewed from $T(c)$, a child J of a critical piece I (say at depth m) is a position on the column zero, say of depth m' , so that marching from J diagonally up right until the depth m , one meets only non-critical positions, except the last one, which is critical.

The critical tableau $T(c)$ is said *persistently recurrent*, if

- every horizontal line contains infinitely many critical positions,
- every vertical line contains at most finitely many critical positions,
- every critical piece has at least one and at most finitely many children.

The key in proving Theorem 1.1 is the following:

Theorem 1.2 *Assume that the critical tableau $T(c)$ is persistently recurrent. Then in the critical nest there are disjoint annuli with moduli μ_n , such that $\mu_n \geq C > 0$ for some C independent of n .*

The implication Theorem 1.2 \implies Theorem 1.1 is quite classical. For details see [QY].

To start with, we recall the three tableaux rules:

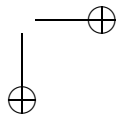
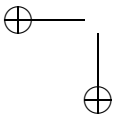
Rule 1 (vertical segment): in each column either there is no critical position or the critical positions form a single vertical segment on the top of the column.

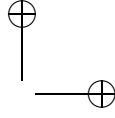
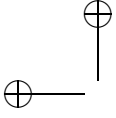
An **upper triangle** in a tableau is by definition a triangle with a vertical side on the left, a horizontal side on top and a diagonal segment as the third side. Its **size** is simply the length of any of its sides, and its **depth** is the depth of its lowest vertex.

Rule 2 (double triangle rule): Two upper triangles in $T(c)$ of the same size and depth such that both vertical sides are critical, are identical.

A **puzzle parallelogram** is a parallelogram with two vertical sides and with the top and lower sides made by diagonal segments.

Rule 3 (double parallelogram rule). Given a puzzle parallelogram D in $T(c)$ whose two vertical sides are critical and whose other pieces are non-





critical, then for any other puzzle parallelogram D' of the same size and depth, with the two top vertices being critical, then either the lower side of D' is either completely non critical, or D' is identical to D .

Lemma 1.3 *Every critical puzzle piece I has at least two children.*

Proof: This is similar to Lemma 1.3 (b) in [M]. Since the convention there is slightly different from ours (we look at the pieces rather than the annuli), we include a proof here.

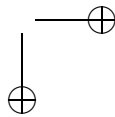
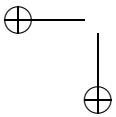
Start from a piece I on the column zero of $T(c)$, say of depth d . March to the right until we first meet a critical position, say at column $k \geq 1$. Denote this position by I' (we know that I and I' represent the same critical piece). Now marching diagonally south-west (left-down) to reach a position S on the column zero. By Rule 1, every position on the open diagonal $diag]S, I'[$ is non critical. Therefore S is a child of I , and is in fact the first child.

By hypothesis of persistent recurrence of $T(c)$ and by Rule 1, the top of the k th-column consists of a segment of finitely critical positions (containing I') and then becomes non-critical from some depth, say d' (with $d' > d$). Denote by P this first non-critical position.

Start now from P . Following first the left down diagonal to reach a piece L on the column zero. By Rule 1 every position on $diag]L, P[$ is non-critical. Now start from P again and proceed up right along the diagonal until the depth of I . There one hits a position, denoted by W . Consider the closed diagonal segment $diag[P, W]$. Apply repeatedly Rule 2 and Rule 3 to each block of k steps starting from P (by comparing with a left most triangle/parallelogram), one concludes that the entire diagonal from P to W , except possibly the last block (of length $< k$), is non-critical. Using then Rule 2 (by comparing with a left most triangle) and then Rule 1 to conclude that every position on the last block, except possibly W , is non-critical. Therefore $diag]L, W[$ is entirely non-critical.

If W is critical then it represents the critical piece I and L is a second child of I . If not march right from W until the first hit of a critical position I'' (representing again the critical piece I). This position exists by assumption. Then by Rule 1 I'' leads to a second child Z of I by following its left down diagonal until the 0th column.

Lemma 1.4 *For any two nested pieces $L \subset\subset K$ in the 0-column of $T(c)$, then $(K \setminus L) \cap P_f = \emptyset$ iff the horizontal strip in $T(c)$ from K to L does not contain semi-critical positions, i.e. no critical vertical segment ends at a depth m with $depth(L) > m \geq depth(K)$.*



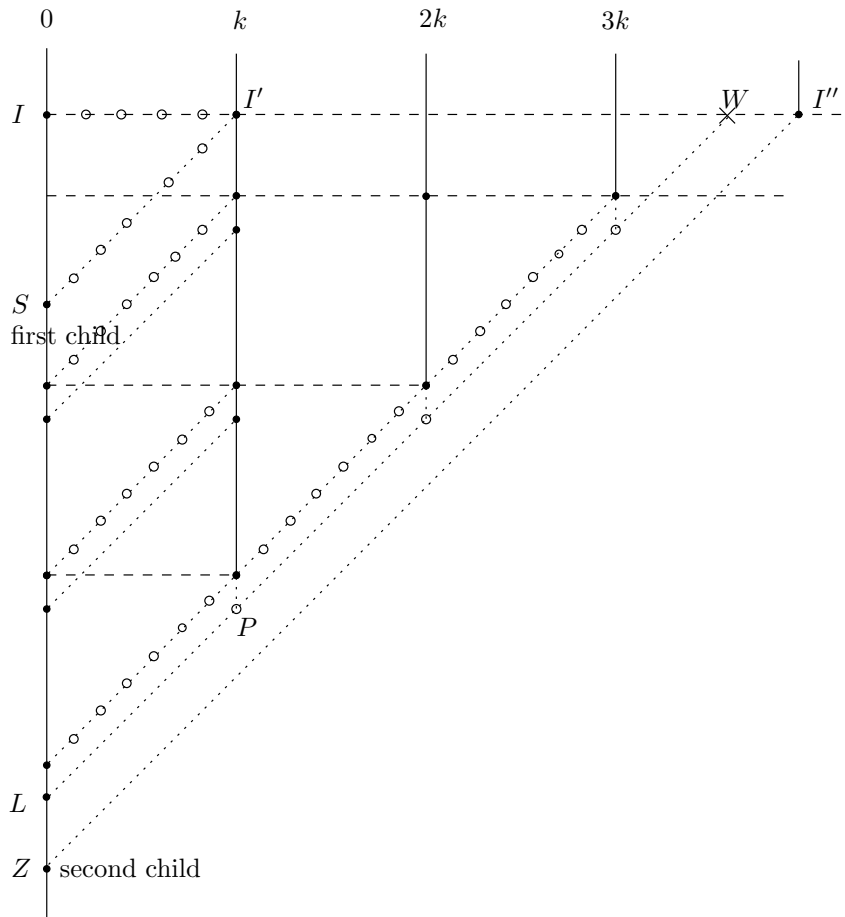
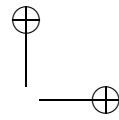
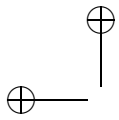


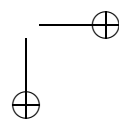
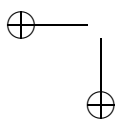
Figure 2.1. The existence of two children (or more).

3 KSS nest

Definition. Let I_0 be a critical puzzle piece. We will define inductively a nested sequence of critical puzzle pieces

$$I_0 \supseteq \cdots \supseteq I_n \supseteq L_n \supseteq K'_n \supseteq K_n \supseteq I_{n+1} \supseteq \cdots,$$

all descendants of I_0 , as follows: Assume I_n is already defined. Then its first and last child are named respectively L_n, K'_n . The piece K_n is the critical puzzle piece contained in K'_n such that the depth difference $|K_n - K'_n|$ is equal to $|L_n - I_n|$. Finally I_{n+1} is the last child of K_n . By Lemma 1.3 all



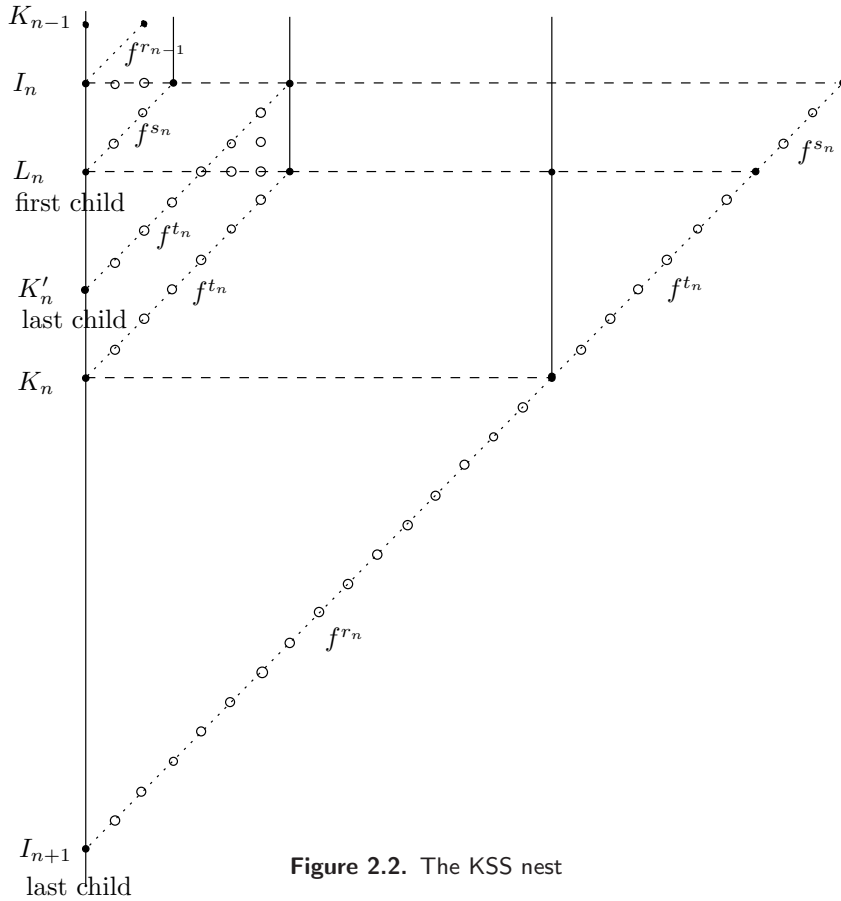
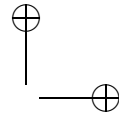
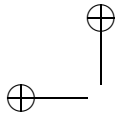
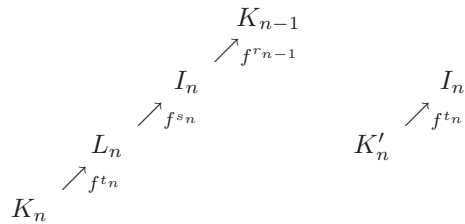


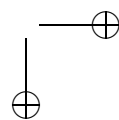
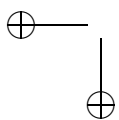
Figure 2.2. The KSS nest

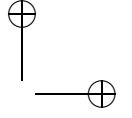
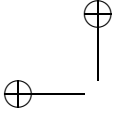
these pieces exist and are mutually distinct.



This nested sequence of puzzle pieces is called *KSS nest* (Kozlovski-Shen-Strien nest). See [KSS] the its original construction.

Here are some basic combinatorial properties of the KSS nest. Set $s_n := |L_n - I_n| = |K_n - K'_n|$, $t_n := |K'_n - I_n|$, $r_n := |I_{n+1} - K_n|$, $p_n := |K_n - K_{n-1}|$.





Lemma 1.5 *We have $f^{r_{n-1}}(I_n) = K_{n-1}$, $f^{s_n}(L_n) = I_n$, $f^{t_n}(K'_n) = I_n$, $f^{t_n}(K_n) = L_n$, each is of degree δ , where $\delta = \deg(f)$, and $f^{p_n}(K_n) = K_{n-1}$, $\deg(f^{p_n} : K_n \rightarrow K_{n-1}) = \delta^3$.*

Proof: The first equalities $f^{r_{n-1}}(I_n) = K_{n-1}$, $f^{s_n}(L_n) = I_n$, $f^{t_n}(K'_n) = I_n$ are almost trivial from the construction: a child S of I is in particular a pullback of I . And the number of pullback iterates is equal to the depth difference.

However the equality $f^{t_n}(K_n) = L_n$ is a lot less trivial. We know that $f^{t_n}(K_n)$ is a piece at the same depth as L_n .

Looking in the tableau $T(c)$, by definition of a child, there is no critical position on the diagonal between K'_n and $f^{t_n}(K'_n)$, except the ends. By Rule 1 there is no critical position on the diagonal between K_n and $f^{t_n}(K_n)$, except probably the north-east end.

Assume by contradiction that $f^{t_n}(K_n)$ is not critical.

Starting from $f^{t_n}(K_n)$ and go to the west (left) until one reaches the diagonal between K'_n and $f^{t_n}(K'_n)$, this horizontal segment has no critical position by Rule 1 and the fact that K'_n is a child.

One goes now from $f^{t_n}(K_n)$ to the east (right) until one hits for the first time a critical position and then turn to going north until the depth of I_n , to reach a piece, which we name I' . Now I' must be critical by Rule 1, i.e. is equal to I_n . Now follow the south-west diagonal from I' until the 0th column, we will find another child of I_n : by Rule 2 and the fact that L_n is the first child the pieces of the diagonal from I' until the depth above that of L_n has no critical positions. When it reaches the depth of L_n it is again non-critical. And then it continues to avoid critical positions all the way to the 0th-column due to Rule 1.

This contradicts the choice of K'_n as the last child. Consequently $f^{t_n}(K_n) = L_n$.

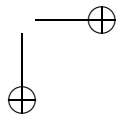
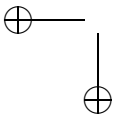
Finally the equality $f^{p_n}(K_n) = K_{n-1}$, i.e. K_n is a pullback of K_{n-1} follows: as a child I_n is a pullback of K_{n-1} , and L_n is a pullback of I_n and finally K_n is a pullback of L_n by above.

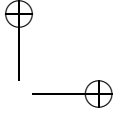
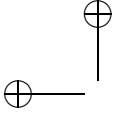
The degree estimates are simple consequences.

The following properties are essential:

Lemma 1.6 *For any $n \geq 1$,*

- (a) *we have $t_n > s_n \geq r_{n-1}$, $p_n = t_n + s_n + r_{n-1}$, $(K'_n \setminus K_n) \cap P_f = \emptyset$;*
- (b) *any two consecutive critical positions on the horizontal line of K_n have a width difference $\geq t_n$ and $\leq r_n$;*
- (c) *we have $r_n \geq 2t_n$, $3t_n \geq p_n \geq 2p_{n-1}$;*
- (d) *for any $m > m^-$, and for $M := |K_m - K_{m^-}|$, we have $f^M(K_m) = K_{m^-}$ of degree $\delta^{3(m-m^-)}$, and $M < 2p_m$.*





Proof: (a). The relations $t_n > s_n \geq r_{n-1}$, $p_n = t_n + s_n + r_{n-1}$ are trivial.

To prove $(K'_n \setminus K_n) \cap P_f = \emptyset$, we apply Lemma 1.4. Assume by contradiction that there is a semi-critical position $Q \supset E$ on the horizontal strip of K'_n and K_n . As Q is critical, we have $Q = K'_n$. As E is not critical, we have $E \neq K_n$. By Rule 2 $f^{t_n}(Q) =: I'$ is critical therefore is equal to I_n , and there is not other critical positions on the diagonal. By Rule 3 $f^{t_n}(E)$ is not critical. Now follow the argument as in the proof of $f^{t_n}(K_n) = L_n$, by going at first to the right of $f^{t_n}(E)$ until the first critical position, then upwards until the depth of I_n and then follow the left-down diagonal until the first critical position. We then find a child of I_n with depth deeper than $\text{depth}(Q) = \text{depth}(K'_n)$, contradicting the fact that K'_n is the last child.

(b) At first we look at the distance between K_n and the first critical position E' on the east (right) of K_n . As the diagonal between K_n and $f^{t_n}(K_n)$ has no critical positions (except the ends), by Rule 1 the width difference $|E' - K_n|$ is at least t_n .

Let now E'', E''' be two consecutive critical positions on the depth of K_n . As E'' is critical it is equal to K_n . By Rule 2 the diagonal between E'' and $f^{t_n}(E'')$ has no critical positions (except the ends), therefore the width difference from E'' to E''' is at least t_n .

Assume that the width difference from E'', E''' is $> r_n$, and that E''' is on the right of E'' . Then starting from E''' and follow the south-west diagonal until we reach a critical position for the first time, we find a child of K_n that is deeper than I_{n+1} , contradicting the choice of I_{n+1} .

(c) Proof of $r_n \geq 2t_n$. As I_{n+1} is a child of K_n , but is not the first child, we have that $f^{r_n}(I_{n+1})$ is some k 'th critical position to the east (right) of K_n with $k \geq 2$. So $r_n \geq 2t_n$. The fact that $p_n = t_n + s_n + r_{n-1} \leq 3t_n$ is due to (a). Now $t_n \geq s_n \geq r_{n-1} \geq 2t_{n-1} \geq 2s_{n-1} \geq 2r_{n-2}$, so $p_n = t_n + s_n + r_{n-1} \geq 2t_{n-1} + 2s_{n-1} + 2r_{n-2} = 2p_{n-1}$.

(d) $M = p_m + p_{m-1} + \dots + p_{m-+1}$. So $M + p_{m-+1} \leq 2p_m$. The rest follows from Lemma 1.5.

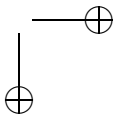
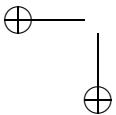
4 From KSS nest to the conditions of KL Lemma

Lemma 1.7 (*Kahn-Lyubich Covering Lemma*) (*[KL]*) Fix $\eta > 0$ and $D \in \mathbb{N}$.

There is $\varepsilon(\eta, D) > 0$ such that:

given any two nests of three hyperbolic discs $A \subset\subset A' \subset\subset U$ and $B \subset\subset B' \subset\subset V$, and any proper holomorphic map $g : U \rightarrow V$ of degree D such that $g|_{A'} : A' \rightarrow B'$ and $g|_A : A \rightarrow B$ are both proper and that $\text{mod}(B' \setminus B) \geq \eta \cdot \text{mod}(U \setminus A)$, then, setting $d := \text{deg}(g|_{A'})$,

$$\text{either } \text{mod}(U \setminus A) > \varepsilon(\eta, D) \quad \text{or} \quad \text{mod}(U \setminus A) > \frac{\eta}{2d^2} \text{mod}(V \setminus B) .$$



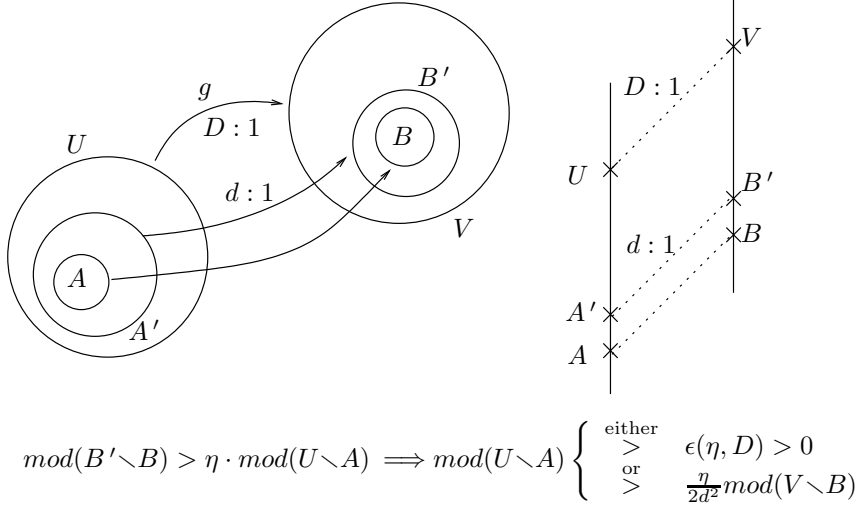


Figure 4.3. Kahn-Lyubich Covering Lemma

This lemma will be our main analytic tool (along side with Grötzsch inequality) to estimate moduli of annuli.

Fix $N > 0$ (which will be $2\delta^{34} + 1$ later on). Fix $m > N$. We will construct a KL-map $g : (U, A', A) \rightarrow (V, B', B)$ depending on the pair (N, m) . Set $U = K_m$ and $V = K_{m-N}$.

Let σ, M be the integers such that $f^\sigma(K'_{m+4}) = K_m = U$, $f^M(K_m) = K_{m-N} = V$. View U, V as to sit on the σ th and the $(\sigma + M)$ th column of $T(c)$. Set $g = f^M|_U : U \rightarrow V$.

Note that $c \in K'_{m+4}$. Set $x = f^\sigma(c)$ and $y = f^M(x)$. As c is persistently recurrent, there is a minimal $s \geq 0$ such that $f^s(y) \in K_m$. This corresponds in $T(c)$ to start from the column of y at $\text{depth}(K_m)$ going to the right $s \geq 0$ steps until the first hit of a critical position. If $s = 0$, i.e. $y \in K_m$, set $B = K_m$ and $B' = K'_m$. If $s > 0$, the $f^{-s}(K_m)$ has a component containing y , denoted by B . As B never visits a critical piece before reaching K_m , the map $f^s : B \rightarrow K_m$ is actually conformal. Now $f^{-s}(K'_m)$ has a component containing y , denoted by B' . As $(K'_m \setminus K_m) \cap P_f = \emptyset$ by Lemma 1.6, we know that $f^s : B' \rightarrow K'_m$ is also conformal. As B, B' are puzzle pieces of depth deeper than that of $K_{m-N} = V$, we have $B \subset\subset B' \subset\subset V$.

Pull back B, B' by $g = f^M$ to get components A, A' containing x .

In order to apply KL-lemma we need to control $D := \text{deg}(g|_U)$ and $d := \text{deg}(g|_{A'})$.

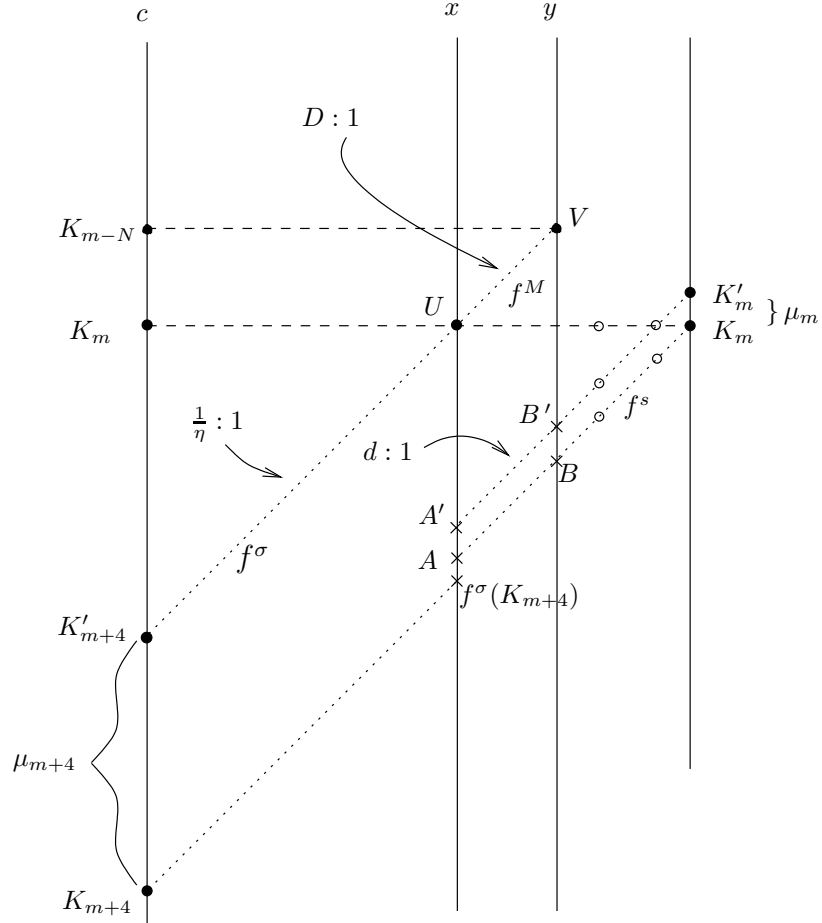
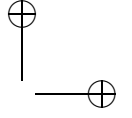
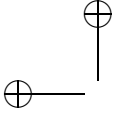


Figure 4.4. Applying Kahn-Lyubich Covering Lemma

Lemma 1.8 Set $\mu_j := \text{mod}(K'_j \setminus K_j)$.

- $\sigma = t_{m+4} + r_{m+3} + p_{m+3} + p_{m+2} + p_{m+1}$ and $\text{deg}(K'_{m+4} \xrightarrow{f^\sigma} K_m) = \delta^{11}$;
- $D := \text{deg}(g : U \rightarrow V) = \delta^{3N}$ (independent of m) and $M < 2p_m$;
- $\#\{0 \leq j < M + s, f^j(A) \subset K_m\} \leq 6$;
- $d := \text{deg}(g : A \rightarrow B) = \text{deg}(g : A' \rightarrow B')$ is at most δ^6 ;
- $f^\sigma(K_{m+4}) \subset A$ and $\text{mod}(U \setminus A) \leq \delta^{11} \cdot \mu_{m+4}$;
- $\text{mod}(V \setminus B) \geq \sum_{j=m-N+1}^m \mu_j$;
- $\text{mod}(B' \setminus B) = \mu_m$.



Proof: a). Note that the number of iterates from K_{m+4} to K_m is $p_{m+4} + p_{m+3} + p_{m+2} + p_{m+1}$. But $|K_{m+4} - K'_{m+4}| = s_{m+4}$. So $\sigma = -s_{m+4} + p_{m+4} + p_{m+3} + p_{m+2} + p_{m+1} = t_{m+4} + r_{m+3} + p_{m+3} + p_{m+2} + p_{m+1}$. Now the total degree is $\delta^{1+1+3+3+3} = \delta^{11}$ by Lemma 1.5.

b). This is due to Lemma 1.6.

c) and d). Note that A is on the σ -th column of $T(c)$, and $f^M(A) = B$. As $f^s(B) = K_m$, we have $\text{depth}(B) - \text{depth}(K_m) = s$ and $\text{depth}(A) - \text{depth}(B) = M$. As $A \subset K_m$, the critical segment on the column of A reaches at least to the depth of K_m .

By Lemma 1.6 two adjacent critical positions on the depth K_m are at least t_m apart. Each visit of $A, f(A), \dots, f^{M-1}(A)$ to K_m corresponds to a critical position. Hence

$$\#\{0 \leq j < M, f^j(A) \subset K_m\} \leq \frac{M}{t_m} + 1 \stackrel{b)}{<} \frac{2p_m}{t_m} + 1 \stackrel{Le.1.6}{\leq} \frac{2 \cdot 3t_m}{t_m} + 1 = 7.$$

Note that $f^M(A)$ (which is B) first visits K_m after s -iterates. So

$$\#\{0 \leq j < M + s, f^j(A) \subset K_m\} = \#\{0 \leq j < M, f^j(A) \subset K_m\} \leq 6.$$

Now for $j \in \{0, 1, \dots, M + s - 1\}$, if $f^j(A)$ is a critical piece then $f^j(A) \subset K_m$. It follows that there are at most six such j and $\text{deg}(f^{M+s} : A \rightarrow K_m) \leq \delta^6$. But $\text{deg}(f^{M+s} : A \rightarrow K_m) = \text{deg}(f^M : A \rightarrow B) = d$. We have $d \leq \delta^6$.

Now B' is a pullback by f^s of K'_m and B is a pullback by f^s of K_m . As $(K'_m \setminus K_m) \cap P_f = \emptyset$, we have that $B' \setminus B$ is a pullback by f^s of $K'_m \setminus K_m$ and is also disjoint from P_f . This implies that $A' \setminus A$ is a pullback by f^M of $B' \setminus B$ and $\text{deg}(f^M|_{A'}) = \text{deg}(f^M|_A)$.

e). Note that $f^\sigma(K_{m+4})$ and A are both puzzle pieces containing x . So one just needs to estimate the depths to see which is contained in which. As they are both pullbacks of K_m we will estimate the number of visits to K_m through iteration:

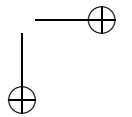
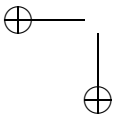
$$K_{m+4} \xrightarrow{f^\sigma} f^\sigma(K_{m+4}) \xrightarrow[a)]{f^{s_{m+4}}} K_m.$$

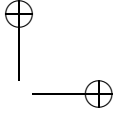
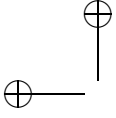
By Lemma 1.6 two adjacent critical positions on the depth K_m are at most r_m apart. So, denoting by $[x]$ the integer part of x ,

$$\#\{0 \leq j < s_{m+4}, f^{j+\sigma}(K_{m+4}) \subset K_m\} \geq \left\lfloor \frac{s_{m+4}}{r_m} \right\rfloor \stackrel{Le.1.6}{\geq} \left\lfloor \frac{r_{m+3}}{r_m} \right\rfloor \stackrel{Le.1.6}{\geq} 8.$$

Therefore $f^\sigma(K_{m+4}) \subset A$. Now

$$\begin{aligned} \text{mod}(U \setminus A) &\stackrel{\text{Grötzsch}}{\leq} \text{mod}(U \setminus f^\sigma(K_{m+4})) = \text{mod}(f^\sigma(K'_{m+4}) \setminus f^\sigma(K_{m+4})) \\ &\stackrel{*}{\leq} \text{deg}(f^\sigma|_{K'_{m+4}}) \cdot \text{mod}(K'_{m+4} \setminus K_{m+4}) = \text{deg}(f^\sigma|_{K'_{m+4}}) \cdot \mu_{m+4} \stackrel{a)}{=} \delta^{11} \mu_{m+4}, \end{aligned}$$





where the inequality marked by $*$ is an equality if f^σ has no critical point on $K'_{m+4} \setminus K_{m+4}$, and a strict inequality otherwise (this can be proved by Grötzsch inequality, see [M], Problem 1-b).

f). Note that B is a piece on the $(\sigma + M)$ th column of $T(c)$. Its critical vertical segment goes down at least to the depth of K_{m-N} . For each $j = m - N + 1, \dots, m$, denote by $l_j \geq 0$ the integer so that $f^{l_j}(y)$ visits K_j for the first time. This corresponds in $T(c)$ to start from the column of y at $\text{depth}(K_j)$ and go to the right $l_j \geq 0$ steps until the first hit of a critical position. We have $l_j \leq s$ by Rule 1. Use f^{l_j} to pullback the annulus $K'_j \setminus K_j$ to get an annulus C_j surrounding B . It is quite easy, using Rule 1 and the fact that $(K'_j \setminus L_j) \cap P_f = \emptyset$, to see that the C_j 's are mutually

disjoint, and $\text{mod}(C_j) = \mu_j$. Therefore $\text{mod}(V \setminus B) \geq \sum_{j=m-N+1}^m \mu_j$.

$$g). \text{mod}(B' \setminus B) = \text{mod}(f^s(B') \setminus f^s(B)) = \text{mod}(K'_m \setminus K_m) = \mu_m.$$

5 Final arguments

Proof of Theorem 1.2. Fix a critical puzzle piece I_0 . Construct the corresponding KSS nest. In particular the annuli $K'_n \setminus K_n$ are mutually disjoint and are nested. Set $\mu_n = \text{mod}(K'_n \setminus K_n)$. We want to show that $\liminf \mu_n > 0$.

Set $N = 2\delta^{34} + 1$. $\forall m > N$, construct $(U, V, \sigma, M, x, y, g, B', B, A', A)$ as above. Then by Lemma 1.8,

$$(*) \begin{cases} \eta := \frac{1}{\text{deg}(f^\sigma|_{K'_{m+4}})} = \frac{1}{\delta^{11}}, \\ D := \text{deg}(f^M|_U) = \delta^{3N}, \\ d := \text{deg}(f^M|_{A'}) \leq \delta^6 \quad \text{and} \quad f^\sigma(K_{m+4}) \subset A. \end{cases}$$

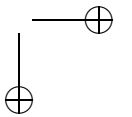
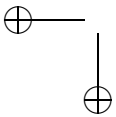
Assume there is a sequence $k_n \rightarrow \infty$ such that $\mu_{m'} \stackrel{m' < k_n}{\geq} \mu_{k_n}$ (otherwise $\liminf \mu_n = \mu_{n_0} > 0$). Fix n such that $k_n - 4 > N$. Set $m = k_n - 4$. Then $\mu_{m'} \geq \mu_{m+4}$ for any $m' \leq m + 4$. Hence,

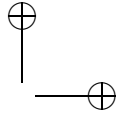
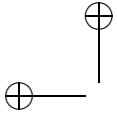
$$\text{mod}(B' \setminus B) \stackrel{g)}{=} \mu_m \stackrel{\text{choice of } m}{\geq} \mu_{m+4} \geq \frac{\text{mod}(U \setminus f^\sigma(K_{m+4}))}{\text{deg}(f^\sigma)} \stackrel{e)}{\geq} \eta \cdot \text{mod}(U \setminus A).$$

We can then apply Kahn-Lyubich Lemma to $g = f^M: (U, A', A) \rightarrow (V, B', B)$, to obtain:

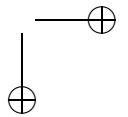
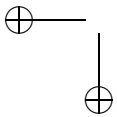
$$\frac{1}{\eta} \mu_{m+4} \geq \text{mod}(U \setminus A) \begin{cases} \text{either} & \varepsilon(\eta, D) > 0; \\ \text{or} & \frac{\eta}{2d^2} \text{mod}(V \setminus B) \stackrel{f)}{\geq} \frac{\eta}{2d^2} (\mu_m + \dots + \mu_{m-N+1}). \end{cases}$$

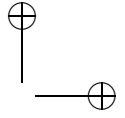
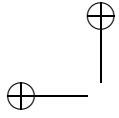
But $\frac{\eta}{2d^2} (\mu_m + \dots + \mu_{m-N+1}) \stackrel{\text{choice of } m}{\geq} \frac{\eta N}{2d^2} \mu_{m+4}$. As $N = 2\delta^{34} + 1 > \frac{2d^2}{\eta^2}$, we get $\mu_{m+4} > \eta \cdot \varepsilon(\eta, D) > 0$.





Therefore $\liminf \mu_n = \lim \mu_{k_n} \geq \eta \cdot \varepsilon(\eta, D) > 0$. (As one can see from the proof, the actual constants in (*) are not important, one needs the following: D is independent of m , the constant η is independent of m, N , and the degree d is bounded by a number independent of m, N .)





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